

# **Fastest Table Sort in the West - Redesigning DuckDB's Sort**

Laurens Kuiper

# Sorting

- One of the most well-studied problems in CS
  - Cache-efficiency
  - Worst-case analysis
  - Parallelism
  - ...
  - ... Arrays!

# Sorting Relational Data

- <sup>D</sup> Use cases:
  - <sup>D</sup> ORDER BY
  - <sup>D</sup> WINDOW
  - <sup>D</sup> Sort-Merge Join
  - <sup>D</sup> Inequality Join
  - <sup>D</sup> ...

# Sorting Relational Data

- Performance challenges
  - Multiple order clauses
  - Different types
  - Columns vs. Rows
  - ...

# Sorting Relational Data

- We distinguish two column types:

```
SELECT *
FROM customer
ORDER BY
  c_birth_country DESC,
  c_birth_year ASC;
```



Payload columns

Key columns

- Both key and payload columns must be ordered!

# The Cost of Sorting

- <sup>D</sup> Dominated by
  - <sup>D</sup> Comparing values
  - <sup>D</sup> Moving data
- <sup>D</sup> This presentation: Comparing key column values
  - <sup>D</sup> Focus: Columnar representation

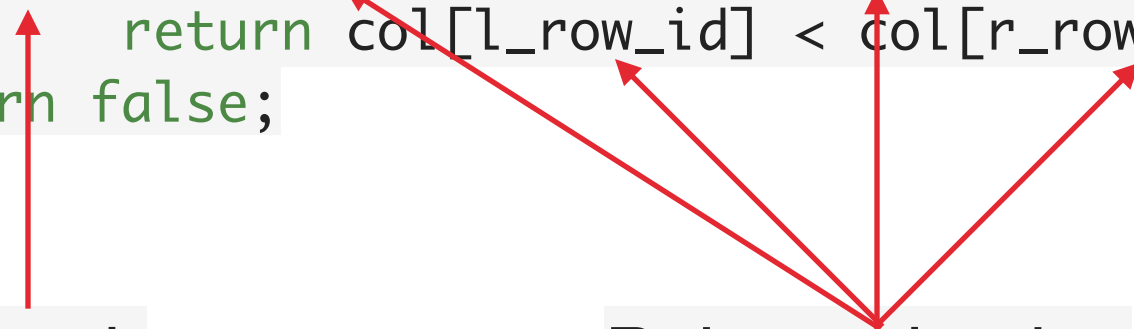
# Comparing Values

- Sorting relational data: row-wise operation
- How to implement comparison for columns?
- Need to sort indices:

```
bool Compare(columns, l_row_id, r_row_id):
    for col : columns:
        if col[l_row_id] != col[r_row_id]:
            return col[l_row_id] < col[r_row_id];
    return false;
```

Branch

Pointer chasing



# Pointer Chasing

- <sup>D</sup> Solution: pack key columns into a row



- <sup>D</sup> No chasing pointers by index
  - <sup>D</sup> Sort data directly
- <sup>D</sup> Better locality
  - <sup>D</sup> Values in row are co-located with each other

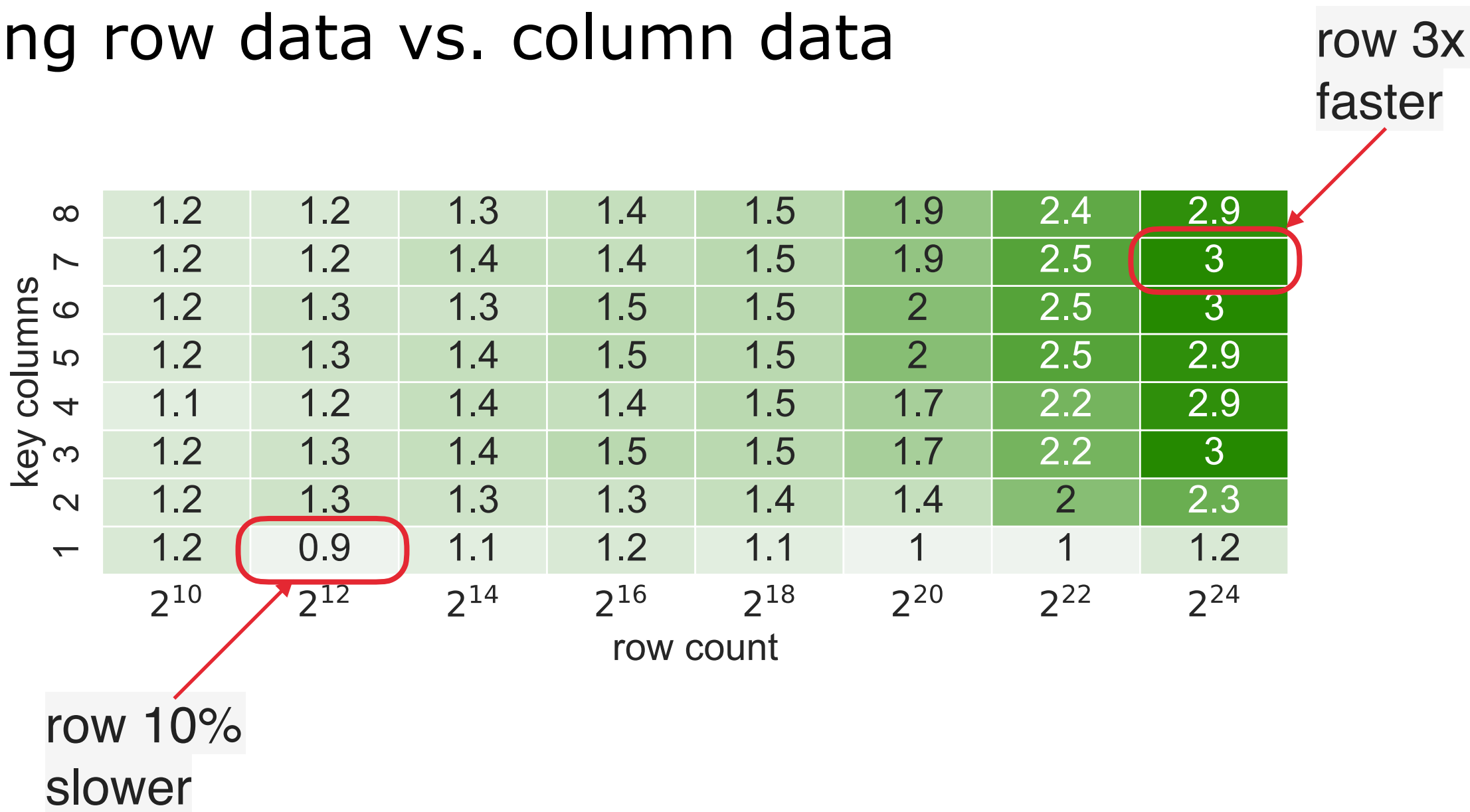


# Pointer Chasing

- <sup>D</sup> Simulated experiment:
  - <sup>D</sup>  $2^{10}$  to  $2^{24}$  tuples
  - <sup>D</sup> 1 to 8 key columns (uint32\_t)
  - <sup>D</sup> Data distribution: many ties
- <sup>D</sup> Measure relative runtime difference
- <sup>D</sup> Hardware: 2020 MacBook Pro (M1)

## Pointer Chasing

- Sorting row data vs. column data



- Time includes column to row transformation!

# Branch Prediction

```
SELECT *
FROM customer
ORDER BY
  c_birth_country DESC,
  c_birth_year ASC;
```

- <sup>D</sup> Comparison has many branches:
  - <sup>D</sup> If c\_birth\_country equal, compare c\_birth\_year
  - <sup>D</sup> ASC/DESC
  - <sup>D</sup> NULLS FIRST/LAST

## Key Normalization

- Encode keys as binary string

```
SELECT *
FROM customer
ORDER BY
  c_birth_country DESC,
  c_birth_year ASC;
```

(a)

birth_country	birth_year
NETHERLANDS	1992
GERMANY	1924

(b)

birth_country	birth_year
78 69 84 72 69 82 76 65 78 68 83 0	200 7 0 0
71 69 82 77 65 78 89 0	132 7 0 0

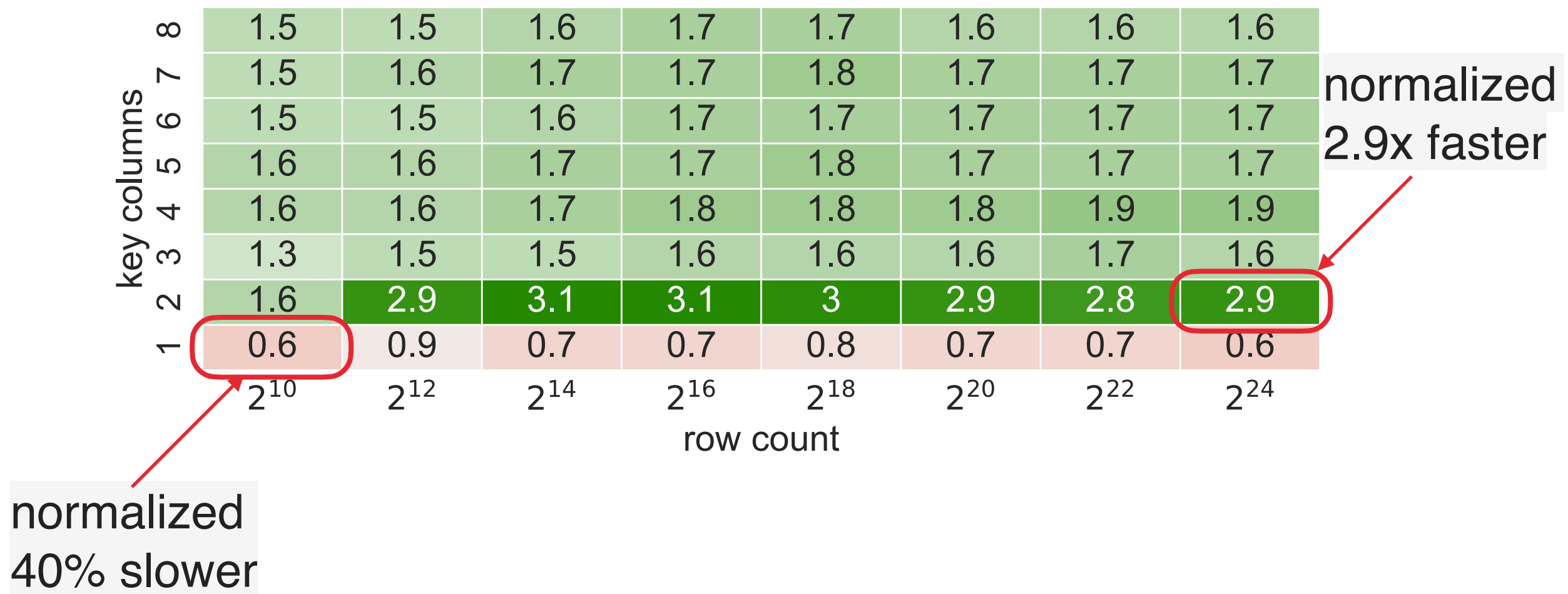
(c)

binary string
177 186 171 183 186 173 179 190 177 187 172 255 128 0 7 200
184 186 173 178 190 177 166 255 255 255 255 255 128 0 7 132

- Eliminates branch predictions from comparison

## Key Normalization

- Sorting normalized row data vs. row data



- Time includes creating the normalized keys!

## Comparing Values: Summary

- Comparing rows: Eliminates pointer chasing
- Key normalization: Eliminates branch predictions
- Both optimizations are almost always worth it

# Sorting Algorithm

- <sup>D</sup> Key normalization enables byte-by-byte Radix Sort
- <sup>D</sup> Key size  $\leq 4$  bytes: LSD
- <sup>D</sup> Key size  $> 4$  bytes: MSD
  - <sup>D</sup> Insertion sort

# Radix Sort vs. Quicksort

- Looks pretty bad for Radix Sort ...

key columns 1 2 3 4 5 6 7 8	0.4	0.3	0.3	0.3	0.3	0.3	0.3	0.3	
	0.5	0.3	0.4	0.4	0.4	0.4	0.4	0.4	
	0.7	0.5	0.8	0.9	0.9	1	1	1	
	0.5	0.5	0.6	0.5	0.6	0.7	0.8	0.7	
	0.6	0.5	0.7	0.6	0.7	0.8	0.7	0.8	
	0.5	0.5	0.6	0.5	0.6	0.7	0.6	0.7	
	0.6	0.5	0.6	0.6	0.7	0.8	0.7	0.8	
	0.5	0.5	0.6	0.5	0.6	0.7	0.6	0.7	
		$2^{10}$	$2^{12}$	$2^{14}$	$2^{16}$	$2^{18}$	$2^{20}$	$2^{22}$	$2^{24}$
		row count							

- ... or does it?



# Radix Sort vs. Quicksort

- <sup>D</sup> Simulation knows key size at compile time
  - <sup>D</sup> In practice we don't
  
- <sup>D</sup> How much does pdqsort benefit from this?

# Radix Sort vs. Quicksort

- Now with dynamic comparison:

key columns	8	0.6	0.5	0.7	0.6	0.7	0.8	0.8	0.9
	7	0.7	0.5	0.7	0.6	0.7	0.9	0.8	1
	6	0.6	0.5	0.6	0.5	0.7	0.8	0.7	0.9
	5	0.7	0.5	0.7	0.6	0.8	1	0.9	1
	4	0.6	0.5	0.7	0.5	0.7	0.9	0.9	0.8
	3	1.7	1.3	2.2	2.5	2.4	2.8	2.9	2.8
	2	1.7	1.5	2	2.2	2.2	2.2	2.2	2.1
	1	1.5	1.5	1.4	1.5	1.4	1.3	1.4	1.4
		$2^{10}$	$2^{12}$	$2^{14}$	$2^{16}$	$2^{18}$	$2^{20}$	$2^{22}$	$2^{24}$
		row count							

- Still cumbersome: many struct definitions

# Radix Sort vs. Quicksort

- <sup>D</sup> To make quicksort efficient we could:
  - <sup>D</sup> Create a lot of templated structs/functions
  - <sup>D</sup> ... and blow up our binary size :(
  - <sup>D</sup> At this point it becomes an arms race



## Sorting Algorithm: Summary

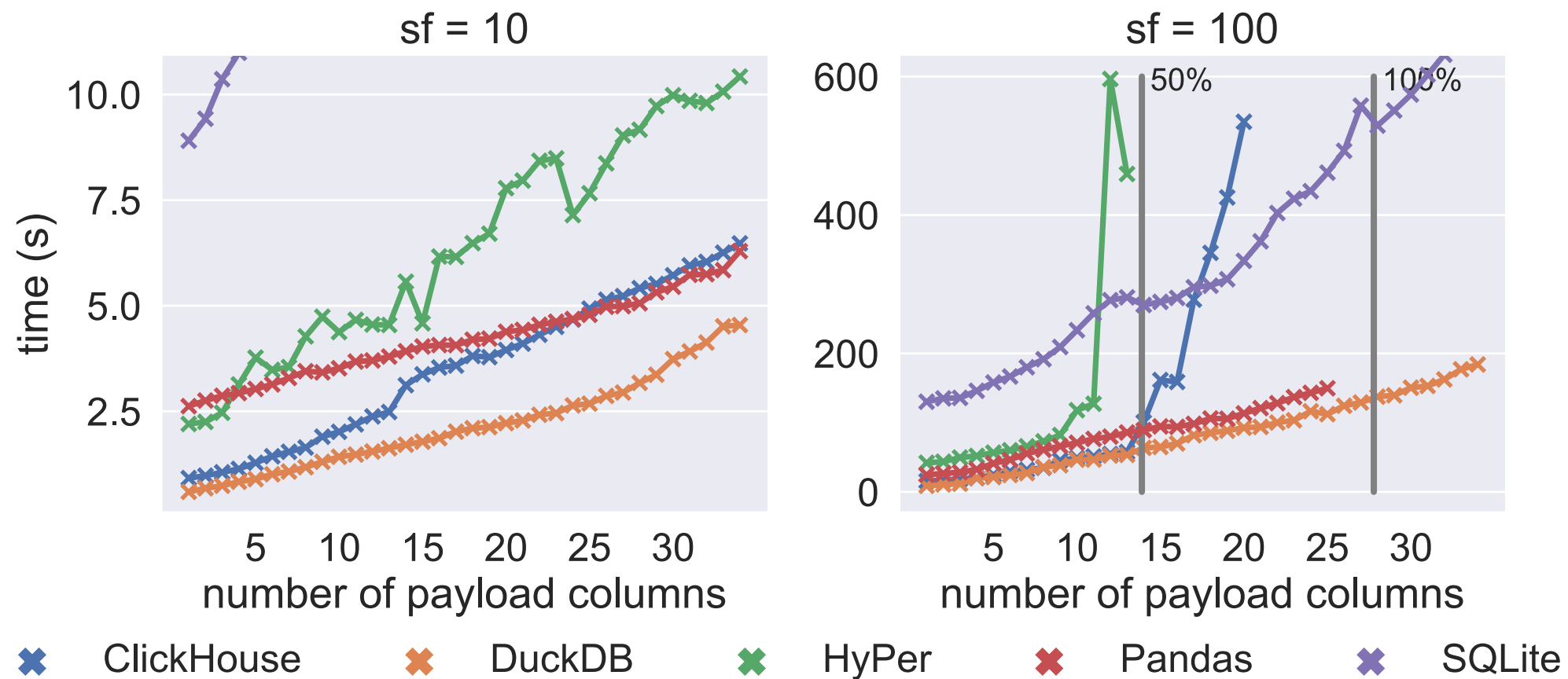
- <sup>D</sup> Performance depends on key size / distribution
- <sup>D</sup> Quicksort needs compile-time optimization
  - <sup>D</sup> Radix Sort does not
- <sup>D</sup> Efficient quicksort is possible for relational data
  - <sup>D</sup> Comes at a cost
  - <sup>D</sup> ... or for free for JIT systems

## End to end performance?

- <sup>D</sup> Relational sorting benchmark using TPC-DS
  - <sup>D</sup> catalog\_sales
    - <sup>D</sup> 34 columns
    - <sup>D</sup> SF10: 14.4M rows
    - <sup>D</sup> SF100: 144M rows
  - <sup>D</sup> Ordered by cs\_quantity, cs\_item\_sk

# TPC-DS catalog\_sales

- Increasing number of payload columns



- Note: 3GB/s SSD write speed

## Wrapping Up

- <sup>D</sup> Sorting relational data efficiently is challenging
- <sup>D</sup> Performance is impacted by:
  - <sup>D</sup> Random access
  - <sup>D</sup> Branch predictions
- <sup>D</sup> We can mitigate these problems with:
  - <sup>D</sup> Row layout in memory
  - <sup>D</sup> Key normalization
- <sup>D</sup> Trade-off between Radix Sort and Quicksort

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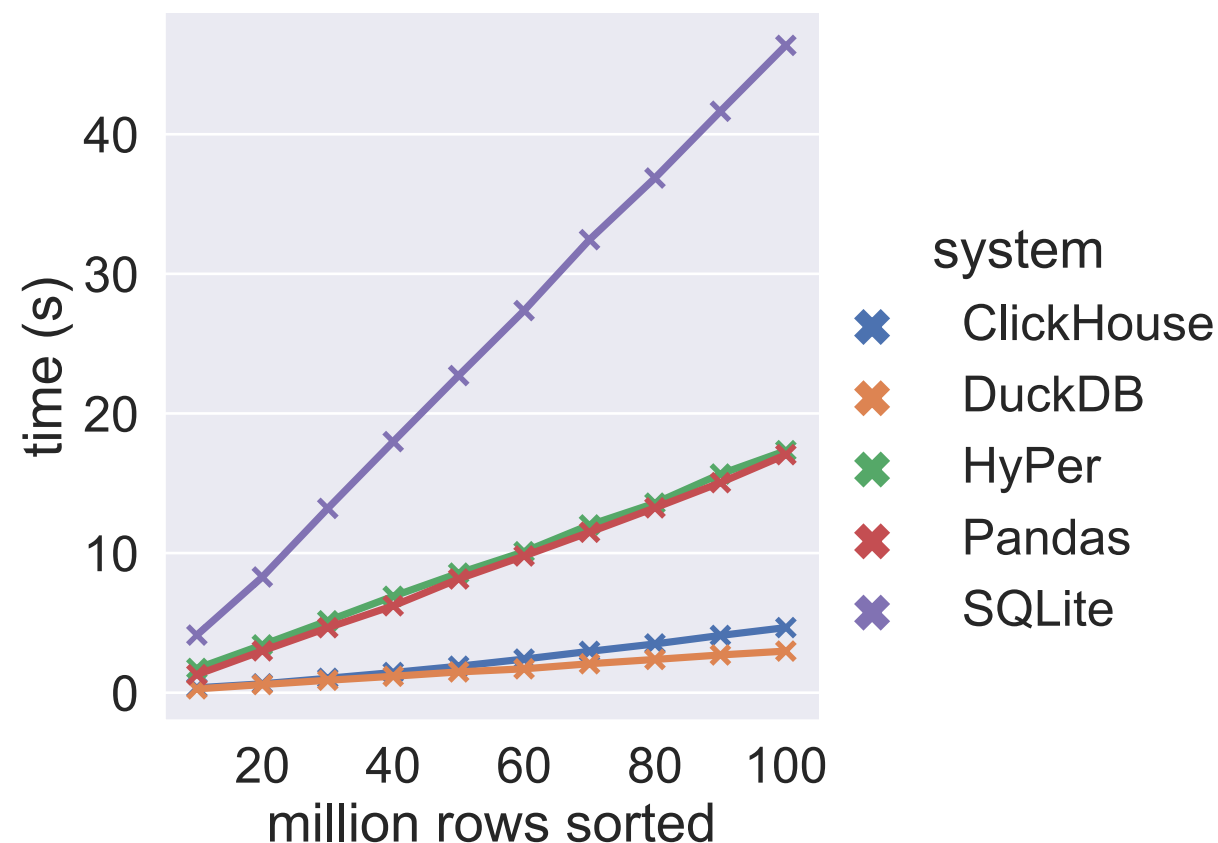


# Key Normalization

- <sup>D</sup> Strings:
  - <sup>D</sup> Encode a prefix
  - <sup>D</sup> Collation can be encoded
- <sup>D</sup> NULL values:
  - <sup>D</sup> Encode using additional byte

## Random Integers

- 10-100M random integers



- ... Actual relational data in the next experiment!

# Parallelism

- <sup>D</sup> DuckDB uses Morsel-Driven Parallelism
  - <sup>D</sup> Threads collect data locally
  - <sup>D</sup> Each thread sorts its own data
- <sup>D</sup> Merge Sort needed for final result!

## Merge Sort

- Parallelized using Merge Path

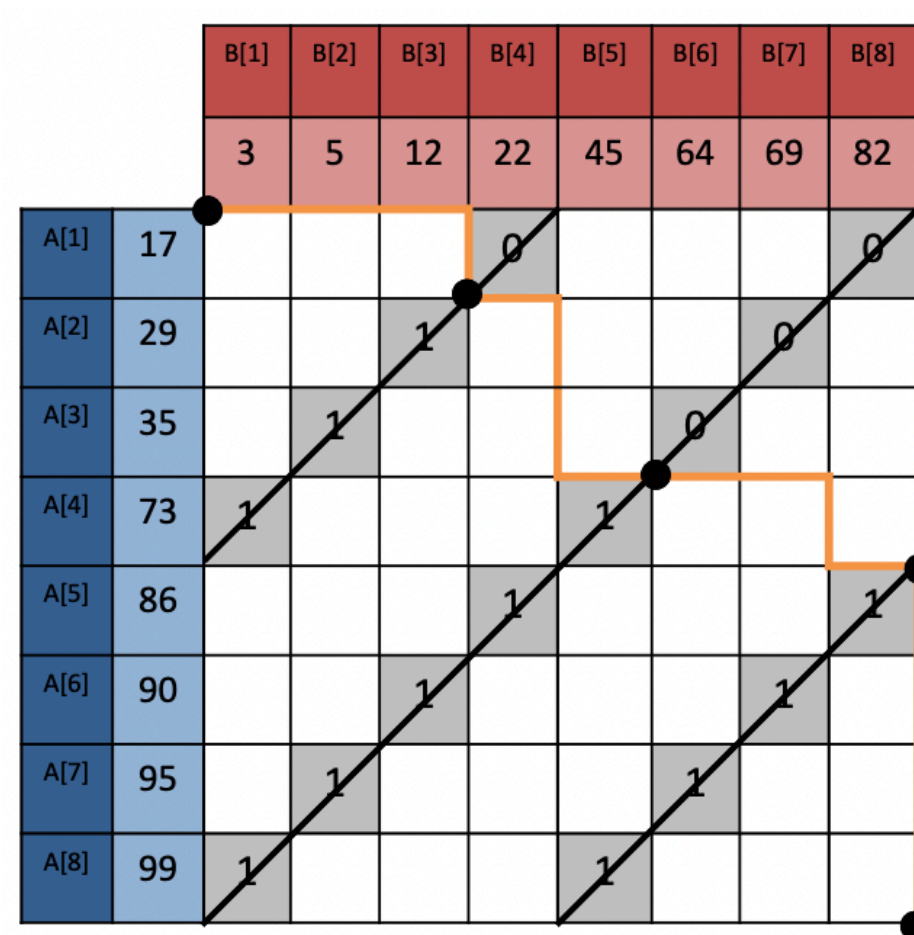
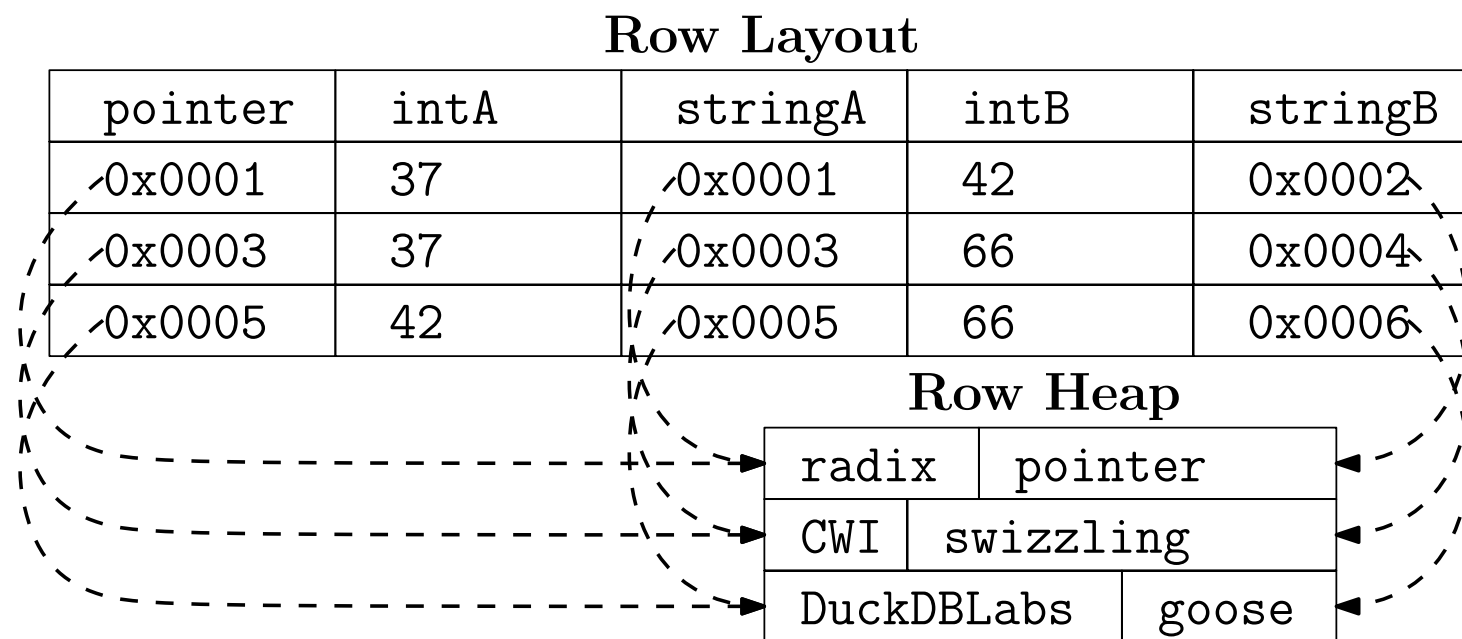


Image by Oded Green et al.

## External Sorting

- Serialize payload in row format
- Buffer-managed blocks



## External Sorting

- Heap also uses rows
- Swizzle pointer -> offset

Row Layout

offset	intA	stringA	intB	stringB
0	37	0	42	5
12	37	0	66	3
24	42	0	66	10

Row Heap

radix	pointer
CWI	swizzling
DuckDBLabs	goose

# External Sorting

- External sorting is made possible because of
  - merge sort
  - buffer manager
  - pointer swizzling
- Modern hardware helps too!